## CECS 528, Homework Assignment 2, Spring 2025, Dr. Ebert

Directions: Please review the Homework section on page 6 of the syllabus including a list of all rules and guidelines for writing and submitting solutions.

Due Date: Monday, February 10th as a PDF-file upload to the HW2 Canvas dropbox.

## Problems

- 1. Apply both the FFT and IFFT algorithms for determining the coefficients of the product C(x) of the linear polynomials A(x) = 2x 5 and B(x) = -3x + 4. Do so as follows.
  - (a) Compute  $DFT_4(A)$  using FFT. (5 points)
  - (b) Compute  $DFT_4(B)$  using FFT. (5 points)
  - (c) Compute  $\vec{c} = \text{DFT}_4(C)$  by multiplying corresponding components of  $\text{DFT}_4(A)$  and  $\text{DFT}_4(B)$ . (5 points)
  - (d) Compute  $DFT_4^{-1}(\vec{c})$  using IFFT. (5 points)
  - (e) Verify that  $DFT_4^{-1}(\vec{c})$  gives the correct coefficients of  $A(x) \cdot B(x)$ . (5 points)
  - (f) Why do we need to work with the fourth roots of unity? What roots of unity would we work with if A and B were both quadratic polynomials? Explain. (5 pts)

Note: please show the calculation for each subproblem of each of the three recursion trees (two FFT trees and one IFFT tree). Points will be lost otherwise.

- 2. Suppose that a CPU of some machine M can, on average, process a single machine instruction in  $5 \times 10^{-9}$  seconds. Moreover, an instance of some problem is an  $n \times n$  matrix R, and there are four known algorithms  $(\mathcal{A}_1, \ldots, \mathcal{A}_4)$  for solving a given problem instance. When these four algorithms are programmed and compiled on M, the average number of instructions (as a function of the number n of rows and columns of matrix R) that need to be executed by each program is i) 1500n, ii)  $250n \log n$ , iii)  $45n^2$ , and iv)  $6n^3$ , respectively. What is the largest input matrix that each program can process and solve within a one-hour period? Show all work for each calculation except to ii) which you may solve with the help of software. (20 points)
- 3. An instance of problem L is an array a of n natural numbers and an integer  $t \ge n$ . Describe an efficient algorithm for finding the least integer d for which

$$\sum_{i=0}^{n-1} \left\lceil \frac{a[i]}{d} \right\rceil \le t$$

First describe your algorithm in one or more paragraphs, followed by providing pseudocode. Finish by analyzing the number of steps your algorithm takes as a function of both n and  $\log t$ , the two size parameters for L. Assume that each arithmetic operation requires O(1) steps. (25 points)