

CECS 528, Homework Assignment 3, Spring 2025, Dr. Ebert

Directions: Please review the Homework section on page 6 of the syllabus including a list of all rules and guidelines for writing and submitting solutions.

Due Date: Monday, February 17th as a PDF-file upload to the HW3 Canvas dropbox.

Problems

1. Given an array $a[0 : n - 1]$ of n positive integers, the goal is to determine the subarray $a[i : j]$ for which

$$Q(i, j) = (j - i + 1) \min_{i \leq k \leq j} (a[k])$$

is maximized, where $0 \leq i \leq j \leq n - 1$. For example if $a = 3, 3, 1, 7, 4, 2, 4, 6, 1$, then $a[3 : 7]$ yields the maximum value, since the minimum a -value in this range is $a[5] = 2$, and

$$Q(3, 7) = (7 - 3 + 1)(2) = 10.$$

One algorithm for finding $a[i : j]$ is to consider all n^2 possible combinations of i and j and keep track of the combination that produces the maximum value. But this algorithm has quadratic running time (why?). Instead, provide a greedy algorithm that can determine the optimal subarray $a[i : j]$ in $O(n \log n)$ steps.

- (a) Clearly describe your algorithm in a few paragraphs, followed by pseudocode. Make sure all variables are clearly defined. (15 pts)
 - (b) Prove that your algorithm's running time $O(n \log n)$. (5 pts)
 - (c) Demonstrate how your algorithm works on the array $a = 5, 9, 10, 2, 8, 3, 7, 1, 6, 4$. Do this by providing a table whose rows represent each stage of the algorithm and whose columns keep track of all necessary variables and structures that are needed to illustrate the algorithm. (10 pts)
 - (d) Using one or more paragraphs, prove that your algorithm will always find the optimal subarray. Note: you may or may not find the Replacement method helpful. It depends on your algorithm. (10 pts)
2. A third algorithm for finding an MST in a weighted graph $G = (V, E, c)$ works as follows. Sort the edges of E from highest to lowest weight. For each edge e taken in sorted order, if removing e does *not* disconnect G , then remove e from E :

$$E \leftarrow E - \{e\}.$$

Otherwise, e remains in E . Then after this procedure, what remains in E is an mst for the original graph G . Use the Replacement method to prove the correctness of this algorithm. Hint: for guidance, study the correctness proofs for Kruskal's and Prim's algorithm. ((15 points)